Query Tuning Using Advanced Hints

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PUSH_PRED Hint

- Applicable when doing an outer join to a view
- Normally, outer join predicate is evaluated after view is materialized
- PUSH_PRED forces the outer join predicate between table and view to be pushed into the view

Outer Join to a View without PUSH_PRED Hint

Optimizer steps

- Execute view query independently
- 2. Materialize view results in an internal temporary table
- Outer Join to the view using sort merge or hash join
 Index is not available for joining to temporary table
 Full scan of temporary table required

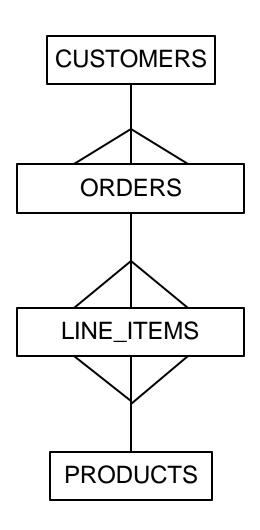
Outer Join to a View with PUSH_PRED Hint

Optimizer steps

- Modify original statement by inserting outer join predicate inside of view statement
- Execute join

If join column in view is indexed then an indexed nested loop join can be performed

PUSH_PRED Hint Sample Tables



PUSH_PRED Hint Sample Tables

CREATE TABLE products (
product_id NUMBER(9),
product_name VARCHAR2(30),
product_descr VARCHAR2(4000),
instock VARCHAR2(1));

PK(Unique Index): product_id

CREATE TABLE orders (order_id NUMBER(9), customer_id NUMBER(6));

PK: order_id

Non-unique index: order_id,customer_id

CREATE TABLE line_items (order_id NUMBER(9), product_id NUMBER(9), comments VARCHAR2(80))

PK(Unique index): product_id, order_id

CREATE TABLE customers (customer_id NUMBER(6), customer_name VARCHAR2(30));

PK(Unique Index): customer_id

PUSH_PRED Hint Query

```
SELECT p. product_name, l. product_id
FROM
  products P,
  (SELECT li.product_id
  FROM line_items li, orders o
  WHERE li.order_id = o.order_id) l
WHERE p. product_i d = l. product_i d(+)
Query Plan without Hint
   SELECT STATEMENT Optimizer=CHOOSE
     HASH JOIN (OUTER)
       TABLE ACCESS (FULL) OF 'PRODUCTS'
       VI EW
         HASH JOIN
            TABLE ACCESS (FULL) OF 'ORDERS'
            TABLE ACCESS (FULL) OF 'LINE_ITEMS'
```

Query with PUSH_PRED Hint

```
SELECT /*+ PUSH_PRED(l) */ p. product_name, l. product_id
FROM
   products P,
   (SELECT li. product_id
   FROM line_items li, orders o
   WHERE li. order_id = o. order_id) l
WHERE p. product_id = l. product_id(+)
```

Query Plan unchanged despite Hint

```
SELECT STATEMENT Optimizer=CHOOSE

HASH JOIN (OUTER)

TABLE ACCESS (FULL) OF 'PRODUCTS'

VIEW

HASH JOIN

TABLE ACCESS (FULL) OF 'ORDERS'

TABLE ACCESS (FULL) OF 'LINE_ITEMS'
```

Parameters that Control Pushing Predicate into View

- PUSH_PRED Hint doesn't work in 8i
- Fixed in 9i
- In 8i need to use undocumented initialization parameter_PUSH_JOIN_PREDICATE
- Set either in init.ora file or with ALTER SESSION statement
- ALTER SESSION SET "_PUSH_JOIN_PREDICATE" = TRUE

Query with PUSH_PRED Hint

```
ALTER SESSION SET "_PUSH_JOIN_PREDICATE" = TRUE;
SELECT /*+ PUSH_PRED(1) */ p. product_name, l. product_i d
FROM
  products P,
  (SELECT li.product_id
  FROM line_items li, orders o
  WHERE li.order_id = o.order_id) l
WHERE p. product_i d = l. product_i d(+)
Query Plan still not changed
   SELECT STATEMENT Optimizer=CHOOSE
     HASH JOIN (OUTER)
       TABLE ACCESS (FULL) OF 'PRODUCTS'
       VI EW
         HASH JOIN
            TABLE ACCESS (FULL) OF 'ORDERS'
            TABLE ACCESS (FULL) OF 'LINE_ITEMS'
```

Index Statistics Affect Pushing Predicate into View

- Optimizer assumes full table scan is cheaper than using index on LINEITEMS(PRODUCT_ID,ORDER_ID)
- Need to convince optimizer that cost of using index on LINEITEMS(PRODUCT_ID,ORDER_ID) is less than cost of full table scan
- Use initialization parameter:

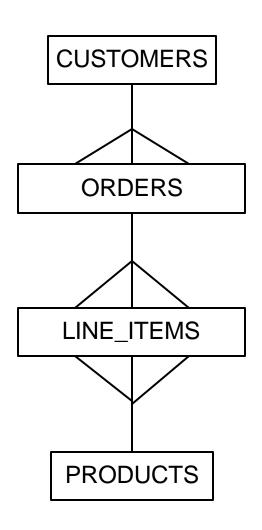
OPTIMIZER_INDEX_COST_ADJ

OPTIMIZER_INDEX_COST_ADJ Initialization Parameter

- Parameter may be set either in init.ora file or using ALTER SESSION statement
- Range of values: 1 .. 10000
- Default value: 100
- Lower values tell optimizer that cost of using an index is lower
- Lower values cause optimizer to favor use of indexes

Query with PUSH_PRED Hint

```
ALTER SESSION SET "_PUSH_JOIN_PREDICATE" = TRUE;
ALTER SESSION SET OPTIMIZER_INDEX_COST_ADJ = 1;
SELECT /*+ PUSH_PRED(1) */ p. product_name, l. product_i d
FROM
 products P.
 (SELECT li.product_id FROM line_items li, orders o
 WHERE li. order id = o. order id) l
WHERE p. product_i d = l. product_i d(+)
Query Plan
   SELECT STATEMENT
      NESTED LOOPS (OUTER)
        TABLE ACCESS (FULL) OF 'PRODUCTS'
        VIEW PUSHED PREDICATE
          NESTED LOOPS
            INDEX (RANGE SCAN) OF 'LINE_ITEMS_PK' (UNIQUE)
            INDEX (RANGE SCAN) OF 'ORDERS_I1' (NON-UNIQUE)
```



```
CREATE TABLE products (
product_id NUMBER(9),
product_name VARCHAR2(30),
product_descr VARCHAR2(4000),
instock VARCHAR2(1));
```

PK(Unique Index): product_id

CREATE TABLE orders (order_id NUMBER(9), customer_id NUMBER(6));

PK: order_id

Non-unique index: order_id,customer_id

CREATE TABLE line_items (order_id NUMBER(9), product_id NUMBER(9), comments VARCHAR2(80))

PK(Unique index): product_id, order_id

CREATE TABLE customers (customer_id NUMBER(6), customer_name VARCHAR2(30));

PK(Unique Index): customer_id

- Query is executed by a currently logged in customer
- Find products where

Product description matches some keyword search criteria

AND

Product is either in stock OR has been bought anytime in the past by the customer

For each matching product:

Show in stock status and whether or not it has been bought by current customer

- Sort matching products with most relevant first
- Show 150 most relevant matching products only
- Use Intermedia Text to do the text search matching

Example query,

Find products with product description containing text, "SUN"

Sample Output

PRODUCT NAME	SCORE	IN STOCK	BOUGHT
Product1	9	Y	N
Product2	9	Y	Y
Product3	9	N	Y
Product4	8	Y	Y

Sample Tables

Number of Rows in Sample Tables

PRODUCTS - 20000

ORDERS - 58000

LINE_ITEMS - 58000

Sample Query: First Attempt

```
SELECT product name, scor, instock, purchased
FROM
  SELECT /*+ ORDERED USE NL(li, o) */
    SCORE(10) scor, product name, instock,
    DECODE(o. order_i d, NULL, 'N', 'Y') purchased
  FROM products p,
    line_items li,
    orders o
  WHERE
    CONTAINS(product_descr, 'SUN', 10) > 0
      AND p. product_i d = li.product_i d(+)
      AND li. order_id = o. order_id(+)
      AND o. customer_i d(+) = 999
      AND (p. instock = 'Y' OR o. order id IS NOT NULL)
  GROUP BY p. product_name, SCORE(10), li.product_id,
    instock
  ORDER BY SCORE(10) DESC
WHERE ROWNUM < 151
```

Sample Query: First Attempt Query Plan/Stats

call	count	cpu	el apsed	di sk	query	current	rows
Parse	1	0. 09	0. 09	0	94	0	0
Execute	1	0. 02	0. 02	0	0	0	0
Fetch	11	3.00	3. 03	0	28153	0	150
total	13	3. 11	3. 14	0	28247	0	150

```
SELECT STATEMENT GOAL: CHOOSE
       COUNT (STOPKEY)
 150
 150
         VI EW
          SORT (GROUP BY STOPKEY)
 150
19702
           FILTER
19702
            NESTED LOOPS (OUTER)
             NESTED LOOPS (OUTER)
19703
              TABLE ACCESS (BY INDEX ROWID) OF 'PRODUCTS'
 201
 201
               DOMAIN INDEX OF 'PRODUCT TEXT'
              INDEX (RANGE SCAN) OF 'LINE_ITEMS_PK' (UNIQUE)
19702
19500
             INDEX (RANGE SCAN) OF 'ORDERS I1' (NON-UNIQUE)
```

Sample Query: Second Attempt

- First find 150 most relevant PRODUCTS that will match criteria
- Then join to LINE_ITEMS
- Number of PRODUCTS and LINE_ITEMS rows joined is fewer
- Less rows accessed/ Better performance

```
SELECT /*+ ORDERED USE_NL(li, ord) */ scor, product_name,
 DECODE(ord. order_i d, NULL, ' N' , ' Y' ) purchased, instock
FROM
  SELECT scor, product name, product id, instock
 FROM
    SELECT SCORE(10) scor, product_name, product_id, instock
    FROM products p
    WHERE CONTAINS(product_descr, 'SUN', 10) > 0
      AND (instock = 'Y' OR)
        EXISTS
          ( SELECT /*+ ORDERED USE_NL(ord) */ 1
             FROM line items li, orders ord
             WHERE ord. customer_id = 999 AND li.order_id = ord.order_id
               AND li.product_id = p.product_id AND ROWNUM = 1
          )
    ORDER BY SCORE(10) DESC
    )
 WHERE ROWNUM < 151
 ) p,
 line items li,
 orders ord
WHERE p. product_i d = li. product_i d(+) AND li. order_i d = ord. order_i d(+) AND
   ord. customer_i d(+) = 999
GROUP BY scor, product_name, li.product_id, instock
ORDER BY scor DESC
```

Sample Query: Second Attempt Query Stats

call	count	cpu	el apsed	di sk	query	current	rows
Parse	1	0. 08	0. 09	0	0	0	0
Execute	1	0.01	0. 01	0	0	0	0
Fetch	11	0. 42	0. 56	0	18140	0	150
total	13	0. 51	0. 66	0	18140	0	150

Sample Query: Second Attempt Query Plan

```
SELECT STATEMENT GOAL: CHOOSE
 150
       SORT (GROUP BY)
14752
        NESTED LOOPS (OUTER)
14753
         NESTED LOOPS (OUTER)
 151
          VI EW
 151
           COUNT (STOPKEY)
 150
            VI EW
 150
             SORT (ORDER BY STOPKEY)
 200
              FILTER
 201
               TABLE ACCESS (BY INDEX ROWID) OF 'PRODUCTS'
                DOMAIN INDEX OF 'PRODUCT_TEXT'
 201
               COUNT (STOPKEY)
   0
                NESTED LOOPS
   0
   0
                 INDEX (RANGE SCAN) OF 'LINE ITEMS PK' (UNIQUE)
   0
                 INDEX (RANGE SCAN) OF 'ORDERS I1' (NON-UNIQUE)
14752
          INDEX (RANGE SCAN) OF 'LINE ITEMS PK' (UNIQUE)
         INDEX (RANGE SCAN) OF 'ORDERS I1' (NON-UNIQUE)
14650
```

Sample Query: Final Attempt

- Do not join to LINE_ITEMS to determine if customer has bought product
- Join to PRODUCTS table again and do correlated subquery to see if there exist any LINES_ITEMS for this product and customer
- Reduces the number of LINE_ITEMS rows that are accessed
- Better performance

```
SELECT /*+ PUSH PRED(op) */ scor, product name,
  DECODE(op. product id, NULL, 'N', 'Y') purchased, instock
FROM
  (
 SELECT scor, product name, product id, instock
  FROM
    SELECT /*+ FIRST ROWS */ score(10) scor, product name, product id, instock
    FROM products p
    WHERE CONTAINS(product_descr, 'SUN', 10) > 0
      AND (instock = 'Y' OR
        EXISTS
          ( SELECT /*+ ORDERED USE NL(ord) */ 1
            FROM line items li, orders ord
            WHERE
              ord. customer_id = 999 AND li.order_id = ord.order_id
              AND li.product_id = p.product_id AND rownum = 1
    ORDER BY SCORE(10) DESC
 WHERE ROWNUM < 151
 ) p,
  SELECT product_id FROM products p
 WHERE EXISTS
    (SELECT /*+ ORDERED USE NL(ORD) */ 1
    FROM line items li, orders ord
    WHERE ord. customer id = 999 AND li. order id = ord. order id
      AND li.product id = p.product id AND ROWNUM = 1
    )
 ) op
WHERE p. product_i d = op. product_i d(+)
ORDER BY scor DESC
```

Final Attempt Query Plan/Stats / No Hint

call	count	cpu	el apsed	di sk	query	current	rows
Parse	1	1. 63	1. 67	0	94	0	0
Execute	1	0. 22	0. 22	0	0	0	0
Fetch	11	1. 46	1. 48	0	132607	0	150
total	13	3. 31	3. 37	0	132701	0	150

Final Attempt Query Plan / No Hint

```
O SELECT STATEMENT
                        GOAL: HINT: FIRST ROWS
  150
        SORT (ORDER BY)
         HASH JOIN (OUTER)
  150
  150
          VI EW
  150
           COUNT (STOPKEY)
  150
            VI EW
  150
             FILTER
  150
              TABLE ACCESS (BY INDEX ROWLD) OF 'PRODUCTS'
               DOMAIN INDEX OF 'PRODUCT_TEXT'
  150
              COUNT (STOPKEY)
    0
               NESTED LOOPS
    0
                INDEX (RANGE SCAN) OF 'LINE ITEMS PROD' (NON-UNIQUE)
   0
                INDEX (RANGE SCAN) OF 'ORDERS I1' (NON-UNIQUE)
   0
  200
          VI EW
  200
           INDEX (FULL SCAN) OF 'PRODUCT PK' (UNIQUE)
20000
            COUNT (STOPKEY)
             NESTED LOOPS
  200
56002
              INDEX (RANGE SCAN) OF 'LINE ITEMS PROD' (NON-UNIQUE)
              INDEX (RANGE SCAN) OF 'ORDERS I1' (NON-UNIQUE)
  200
```

Final Attempt Query Plan/Stats / With Hint

call	count	cpu	el apsed	di sk	query	current	rows
Parse	1	0. 01	0. 01	0	0	0	0
Execute	1	0.01	0. 01	0	0	0	0
Fetch	11	0. 07	0. 07	0	1392	0	150
total	13	0. 09	0. 09	0	1392	0	150

Final Attempt Query Plan / With Hint

```
O SELECT STATEMENT
                      GOAL: HINT: FIRST ROWS
150
     SORT (ORDER BY)
      NESTED LOOPS (OUTER)
150
151
       VI EW
151
        COUNT (STOPKEY)
150
         VI EW
150
          FILTER
150
            TABLE ACCESS (BY INDEX ROWLD) OF 'PRODUCTS'
            DOMAIN INDEX OF 'PRODUCT_TEXT'
150
            COUNT (STOPKEY)
 0
             NESTED LOOPS
 0
             INDEX (RANGE SCAN) OF 'LINE ITEMS PK' (UNIQUE)
 0
 0
             INDEX (RANGE SCAN) OF 'ORDERS I1' (NON-UNIQUE)
       VIEW PUSHED PREDICATE
150
150
         INDEX (UNIQUE SCAN) OF 'PRODUCT PK' (UNIQUE)
150
          COUNT (STOPKEY)
           NESTED LOOPS
150
330
           INDEX (RANGE SCAN) OF 'LINE ITEMS PK' (UNIQUE)
           INDEX (RANGE SCAN) OF 'ORDERS I1' (NON-UNIQUE)
150
```

HASH_AJ Hint

- Use for NOT IN queries where there is no index on column in nested query
- Find rows in table A where there are no matching rows in table B
- Find CUSTOMERS who have no ORDERS
- There is no index on ORDERS.CUSTOMER_ID

NOT IN Example / No Hint

SELECT count(*) FROM customers
WHERE customer_id NOT IN
(SELECT customer_id FROM orders)

NOT IN / No Hint Query Stats/Plan

call	count	cpu	el apsed	di sk	query	current	rows
Parse	1	0. 00	0. 00	0	0	0	0
Execute	1	0.00	0.00	0	0	0	0
Fetch	2	67. 16	67. 24	0	141975	4	1
total	4	67. 16	67. 24	0	141975	4	1

```
O SELECT STATEMENT GOAL: CHOOSE
```

1 SORT (AGGREGATE)

998 FILTER

1000 INDEX (FAST FULL SCAN) OF 'CUSTOMER_PK' (UNIQUE)

999 INDEX (FULL SCAN) OF 'ORDERS_I1' (NON-UNIQUE)

NOT IN Example / No Hint

- For each row in CUSTOMERS do a full scan of ORDERS to find a matching row
- 999 CUSTOMERS rows times 58000 ORDERS rows ~= 58,000,000
- ~1000 * 58000 = 58,000,000 row accesses required to execute this query
- This explains high number of block accesses
- Very expensive

NOT IN Example / With HASH_AJ Hint

```
SELECT count(*) FROM customers
WHERE customer_id NOT IN
  (SELECT /*+ HASH_AJ */ customer_id
  FROM orders)
```

NOT IN / With HASH_AJ Hint / Query Plan/Stats

call	count	сри	el apsed	di sk	query	current	rows
Parse	1	0. 01	0. 01	0	0	0	0
Execute	1	0.00	0.00	0	0	0	0
Fetch	2	0. 14	0. 14	0	118	8	1
total	4	0. 15	0. 15	0	118	8	1

```
O SELECT STATEMENT GOAL: CHOOSE
```

1 SORT (AGGREGATE)

998 HASH JOIN (ANTI)

999 INDEX (FAST FULL SCAN) OF 'CUSTOMER_PK' (UNIQUE)

58000 VI EW OF ' VW_NSO_1'

58000 TABLE ACCESS (FULL) OF 'ORDERS'

NOT IN Example / With HASH_AJ Hint

- HASH_AJ Hint forces optimizer to perform a HASH ANTI-JOIN
- HASH ANTI-JOIN is similar to HASH JOIN
- Instead of finding matching rows, find rows in one table with no matching rows in other table
- Each table in HASH ANTI-JOIN is scanned only once
- 999 CUSTOMERS rows + 58000 ORDERS rows = 58999 row accesses
- 58,999 is far less than 58,000,000 row accesses without HASH_AJ hint
- This explains low number of block accesses in stats

HASH_AJ Prerequisites

- Join columns must be NOT NULL
- Assume ORDERS.CUSTOMER_ID or LINE_ITEMS.CUSTOMER_ID were NULLABLE
- Query must be rewritten as

```
SELECT count(*) FROM customers
WHERE customer_id IS NOT NULL AND
  customer_id NOT IN
    (SELECT /*+ HASH_AJ */ customer_id
    FROM orders
  WHERE customer_id IS NOT NULL)
```

NOT EXISTS

■ MYTH:

NOT EXISTS always better than NOT IN

- Assume ORDERS.CUSTOMER_ID is indexed
- 58000 rows in CUSTOMERS table

```
SELECT count(*)
FROM customers c
WHERE NOT EXISTS
   (SELECT 1 FROM
    orders o
WHERE o. customer_id = c. customer_id);
```

NOT EXISTS Versus NOT IN With HASH_AJ

NOT Exists with Index on ORDERS.CUSTOMER_ID

call	count	cpu	el apsed	di sk	query	current	rows
Parse	1	0. 01	0. 01	0	0	0	0
Execute	1	0.00	0.00	0	0	0	0
Fetch	2	1. 63	1. 63	0	116110	4	1
total	4	1. 64	1. 64	0	116110	4	1

NOT IN with HASH_AJ Hint

call	count	cpu	el apsed	di sk	query	current	rows
Parse	1	0. 00	0. 00	0	0	0	0
Execute	1	0.00	0.00	0	0	0	0
Fetch	2	0. 35	0. 36	0	225	8	1
total	4	0. 35	0. 36	0	225	8	1

NOT EXISTS Versus NOT IN with HASH_AJ

- Case 1: No Index on Inner Table
 - NOT IN with HASH_AJ is better
- Case 2: Index on Inner Table
 - Number of rows in Outer Table is Large
 - NOT IN with HASH_AJ is better
 - Example: 58000 rows in CUSTOMERS table
- Case 3: Index on Inner Table
 - Number of rows in Outer Table is Small
 - Number of rows in Inner Table is Large
 - NOT EXISTS is usually better
- Size of Inner Table may also affect result
 - NOT EXISTS may perform better for large indexed inner table

HASH_SJ Hint

- Use for CORRELATED EXISTS queries where there is no index on column in nested query
- Find rows in table A where there exist matching rows in table B
- Find CUSTOMERS who have one or more ORDERS
- There is no index on ORDERS.CUSTOMER_ID

HASH_SJ Example

```
SELECT count(*) FROM customers c
WHERE EXISTS
  (SELECT 1 /*+ HASH_SJ */ FROM orders o
WHERE c. customer_i d = o. customer_i d)
```

HASH_SJ Example Query Stats/Plan

call	count	сри	el apsed	di sk	query	current	rows
Parse	1	0. 02	0. 05	0	0	0	0
Execute	1	0.00	0.00	0	0	0	0
Fetch	2	0. 12	0. 12	0	118	8	1
total	4	0. 14	0. 17	0	118	8	1

- O SELECT STATEMENT GOAL: CHOOSE
- 1 SORT (AGGREGATE)
- 1 HASH JOIN (SEMI)
- 999 INDEX (FAST FULL SCAN) OF 'CUSTOMER_PK' (UNIQUE)
- 58000 TABLE ACCESS (FULL) OF 'ORDERS'

HASH_SJ Example

- HASH_SJ Hint forces optimizer to perform a HASH SEMI-JOIN
- HASH SEMI-JOIN is similar to HASH JOIN
- Unlike true HASH JOIN, only one row from outer table is returned even if there is more than one matching row from inner table
- Each table in HASH SEMI-JOIN is scanned only once
- 999 CUSTOMERS rows + 58000 ORDERS rows = 58999 row accesses
- This are far less than number of row accesses required without HASH_SJ hint
- This explains low number of block accesses in stats

Correlated EXISTS with Index on Inner Table

Example:

```
SELECT count(*) FROM customers c
WHERE EXISTS
  (SELECT 1 FROM orders o
WHERE c. customer_i d = o. customer_i d)
```

- Index on ORDERS.CUSTOMER_ID
- HASH_SJ Hint can yield better performance if number of rows in outer table (CUSTOMERS) is large

HASH_AREA_SIZE Init Parameter

- HASH_AREA_SIZE determines amount of memory allocated for HASH ANTI-JOIN, HASH SEMI-JOIN and HASH JOIN.
- Set HASH_AREA_SIZE appropriately when using HASH_AJ or HASH_SJ Hint on large tables
- Settable via
 - init.ora file
 - ALTER SYSTEM SET HASH_AREA_SIZE=
 - ALTER SESSION SET HASH_AREA_SIZE=

PUSH_SUBQ Hint

- Use to force evaluation of correlated subqueries as early as possible in a query involving joins
- Optimizer usually evaluates subqueries after evaluating joins
- PUSH_SUBQ hint causes optimizer to evaluate subqueries before evaluating all joins

PUSH_SUBQ Hint Example

Table

```
CREATE TABLE shipments (
shipment_id NUMBER(6),
order_id NUMBER(6),
shipment_date DATE);

PK(Unique Index): shipment_id
Index On: order_id
```

Index

Index on line_items(order_id,product_id)

PUSH_SUBQ Example / No Hint

Find all orders and order lines for orders that have shipped

```
SELECT /*+ ORDERED USE_NL(L, P) */ o. order_id,
   l. comments, p. product_name
FROM orders o, line_items l, products p
WHERE o. order_id = l. order_id
   AND l. product_id = p. product_id
   AND EXISTS
   (SELECT 1
   FROM shipments s
   WHERE s. order_id = o. order_id)
AND o. order_id < 50000;</pre>
```

PUSH_SUBQ Example / No Hint: Query Stats

cal l	count	cpu	el apsed	di sk	query	current	rows
Parse	1	0. 00	0. 00	0	0	0	0
Execute	1	0.00	0.00	0	0	0	0
Fetch	1000	5. 03	5. 06	0	417785	4	999
total	1002	5. 03	5. 06	0	417785	4	999

PUSH_SUBQ Example / No Hint: Query Plan

```
SELECT STATEMENT GOAL: CHOOSE
 999
       FI LTER
44502 NESTED LOOPS
47000
         NESTED LOOPS
          INDEX (FAST FULL SCAN) OF 'ORDERS_I1' (NON-UNIQUE)
47000
          TABLE ACCESS (BY INDEX ROWID) OF 'LINE ITEMS'
93998
           INDEX (RANGE SCAN) OF 'LINE ITEMS ORD' (NON-UNIQUE)
93998
91500
         TABLE ACCESS (BY INDEX ROWID) OF 'PRODUCTS'
          INDEX (UNIQUE SCAN) OF 'PRODUCT PK' (UNIQUE)
91500
44501
        INDEX (RANGE SCAN) OF 'SHIPMENTS ORD' (NON-UNIQUE)
```

PUSH_SUBQ Example / No Hint

- Optimizer finds 47000 matching ORDERS rows
- It joins 47000 ORDERS rows with LINE_ITEMS and PRODUCTS
- After join is processed 44502 rows are found
- For each row in the join result the correlated subquery is executed
- Correlated subquery filters out all but 999 rows
- Very expensive

PUSH_SUBQ Example / With Hint

PUSH_SUBQ Hint placed in outer query block

```
SELECT /*+ ORDERED USE_NL(L, P) PUSH_SUBQ */ o. order_id,
   l. comments, p. product_name
FROM orders o, line_items l, products p
WHERE o. order_id = l. order_id
   AND l. product_id = p. product_id
   AND EXISTS
    (SELECT 1
    FROM shipments s
   WHERE s. order_id = o. order_id)
AND o. order_id < 50000;</pre>
```

PUSH_SUBQ / With Hint: Query Stats

call	count	cpu	el apsed	di sk	query	current	rows
Parse	1	0. 00	0. 00	0	0	0	0
Execute	1	0.00	0.00	0	0	0	0
Fetch	1000	1. 39	1. 52	0	103149	4	999
total	1002	1. 39	1. 52	0	103149	4	999

PUSH_SUBQ / With Hint: Query Plan

```
SELECT STATEMENT GOAL: CHOOSE
 999
       NESTED LOOPS
 1000
        NESTED LOOPS
 1000
         INDEX (FAST FULL SCAN) OF 'ORDERS I1' (NON-UNIQUE)
          INDEX (RANGE SCAN) OF 'SHIPMENTS_ORD' (NON-UNIQUE)
46999
          TABLE ACCESS (BY INDEX ROWLD) OF 'LINE_ITEMS'
1998
 1998
          INDEX (RANGE SCAN) OF 'LINE ITEMS ORD' (NON-UNIQUE)
        TABLE ACCESS (BY INDEX ROWID) OF 'PRODUCTS'
 999
 1998
         INDEX (UNIQUE SCAN) OF 'PRODUCT PK' (UNIQUE)
```

PUSH_SUBQ Example / No Hint

- Optimizer first finds matching ORDERS rows
- For each matching ORDERS row it executes the correlated subquery against SHIPMENTS
- 999 ORDERS remain
- 999 ORDERS rows are joined with LINE_ITEMS and PRODUCTS
- Joining 999 ORDERS rows with LINE_ITEMS and PRODUCTS is much more efficient than joining 47000 ORDERS rows
- This explains lower number of block accesses

PUSH_SUBQ Hint

Use When

- There is a correlated subquery that correlates back to a table early in join order
- The correlated subquery substantially reduces the number of rows in the outer table

Benefit

 Filtering tables early in join order prior to doing join substantially reduces the number of rows accessed in tables that occur later in join order